

## Example

Determine the Crout factorization of the symmetric tridiagonal matrix

$$\begin{bmatrix} 2 & -1 & 0 & 0 \\ -1 & 2 & -1 & 0 \\ 0 & -1 & 2 & -1 \\ 0 & 0 & -1 & 2 \end{bmatrix},$$

and use this factorization to solve the linear system

$$\begin{aligned} 2x_1 - x_2 &= 1, \\ -x_1 + 2x_2 - x_3 &= 0, \\ -x_2 + 2x_3 - x_4 &= 0, \\ -x_3 + 2x_4 &= 1. \end{aligned}$$

**Solution** The  $LU$  factorization of  $A$  has the form

$$A = \begin{bmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{34} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & 0 & 0 \\ l_{21} & l_{22} & 0 & 0 \\ 0 & l_{32} & l_{33} & 0 \\ 0 & 0 & l_{43} & l_{44} \end{bmatrix} \begin{bmatrix} 1 & u_{12} & 0 & 0 \\ 0 & 1 & u_{23} & 0 \\ 0 & 0 & 1 & u_{34} \\ 0 & 0 & 0 & 1 \end{bmatrix}$$
$$= \begin{bmatrix} l_{11} & l_{11}u_{12} & 0 & 0 \\ l_{21} & l_{22} + l_{21}u_{12} & l_{22}u_{23} & 0 \\ 0 & l_{32} & l_{33} + l_{32}u_{23} & l_{33}u_{34} \\ 0 & 0 & l_{43} & l_{44} + l_{43}u_{34} \end{bmatrix}$$

Thus

$$\begin{array}{ll} a_{11} : & 2 = l_{11} \implies l_{11} = 2, \\ a_{21} : & -1 = l_{21} \implies l_{21} = -1, \\ a_{23} : & -1 = l_{22}u_{23} \implies u_{23} = -\frac{2}{3}, \\ a_{33} : & 2 = l_{33} + l_{32}u_{23} \implies l_{33} = \frac{4}{3}, \\ a_{43} : & -1 = l_{43} \implies l_{43} = -1, \\ a_{12} : & -1 = l_{11}u_{12} \implies u_{12} = -\frac{1}{2}, \\ a_{22} : & 2 = l_{22} + l_{21}u_{12} \implies l_{22} = -\frac{3}{2}, \\ a_{32} : & -1 = l_{32} \implies l_{32} = -1, \\ a_{34} : & -1 = l_{33}u_{34} \implies u_{34} = -\frac{3}{4}, \\ a_{44} : & 2 = l_{44} + l_{43}u_{34} \implies l_{44} = \frac{5}{4}. \end{array}$$

This gives the Crout factorization

$$A = \begin{bmatrix} 2 & -1 & 0 & 0 \\ -1 & 2 & -1 & 0 \\ 0 & -1 & 2 & -1 \\ 0 & 0 & -1 & 2 \end{bmatrix} = \begin{bmatrix} 2 & 0 & 0 & 0 \\ -1 & \frac{3}{2} & 0 & 0 \\ 0 & -1 & \frac{4}{3} & 0 \\ 0 & 0 & -1 & \frac{5}{4} \end{bmatrix} \begin{bmatrix} 1 & -\frac{1}{2} & 0 & 0 \\ 0 & 1 & -\frac{2}{3} & 0 \\ 0 & 0 & 1 & -\frac{3}{4} \\ 0 & 0 & 0 & 1 \end{bmatrix} = LU.$$

Solving the system

$$Lz = \begin{bmatrix} 2 & 0 & 0 & 0 \\ -1 & \frac{3}{2} & 0 & 0 \\ 0 & -1 & \frac{4}{3} & 0 \\ 0 & 0 & -1 & \frac{5}{4} \end{bmatrix} \begin{bmatrix} z_1 \\ z_2 \\ z_3 \\ z_4 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \end{bmatrix} \quad \text{gives} \quad \begin{bmatrix} z_1 \\ z_2 \\ z_3 \\ z_4 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{3} \\ \frac{1}{4} \\ 1 \end{bmatrix},$$

and then solving

$$Ux = \begin{bmatrix} 1 & -\frac{1}{2} & 0 & 0 \\ 0 & 1 & -\frac{2}{3} & 0 \\ 0 & 0 & 1 & -\frac{3}{4} \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{3} \\ \frac{1}{4} \\ 1 \end{bmatrix} \quad \text{gives} \quad \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}. \quad \blacksquare$$

The Crout Factorization Algorithm can be applied whenever  $l_{ii} \neq 0$

This is true for positive definite matrices and also for strictly diagonally dominant matrices.

### **Theorem**

Suppose that  $A = [a_{ij}]$  is tridiagonal with  $a_{i,i-1}a_{i,i+1} \neq 0$ , for each  $i = 2, 3, \dots, n-1$ . If  $|a_{11}| > |a_{12}|$ ,  $|a_{ii}| \geq |a_{i,i-1}| + |a_{i,i+1}|$ , for each  $i = 2, 3, \dots, n-1$ , and  $|a_{nn}| > |a_{n,n-1}|$ , then  $A$  is nonsingular and the values of  $l_{ii}$  described in the Crout Factorization Algorithm are nonzero for each  $i = 1, 2, \dots, n$ .

### **HOMEWORK 9:**

Exercise Set 6.6: 13,14

# Computational Cost

- The computational cost of LU factorization and Gaussian elimination method is almost the same ( $O(n^3/3)$ ). The LU factorization is efficient for several linear systems with the same coefficient matrix.
- The inverse matrix method is not computationally efficient.
- The  $LDL^t$  or  $LL^t$  factorization for positive definite matrices, reduces the computational cost rather than LU factorization method.
- The computational cost of  $LL^t$  method is less than  $LDL^t$  method.
- The Crout factorization for tridiagonal matrices, reduces computational cost rather than LU factorization method.

# Iterative Techniques for Solving Linear Systems

Matrices with high percentage of zero entries (sparse matrices) are often solved using iterative methods.

## ***Definition***

A **vector norm** on  $\mathbb{R}^n$  is a function,  $\| \cdot \|$ , from  $\mathbb{R}^n$  into  $\mathbb{R}$  with the following properties:

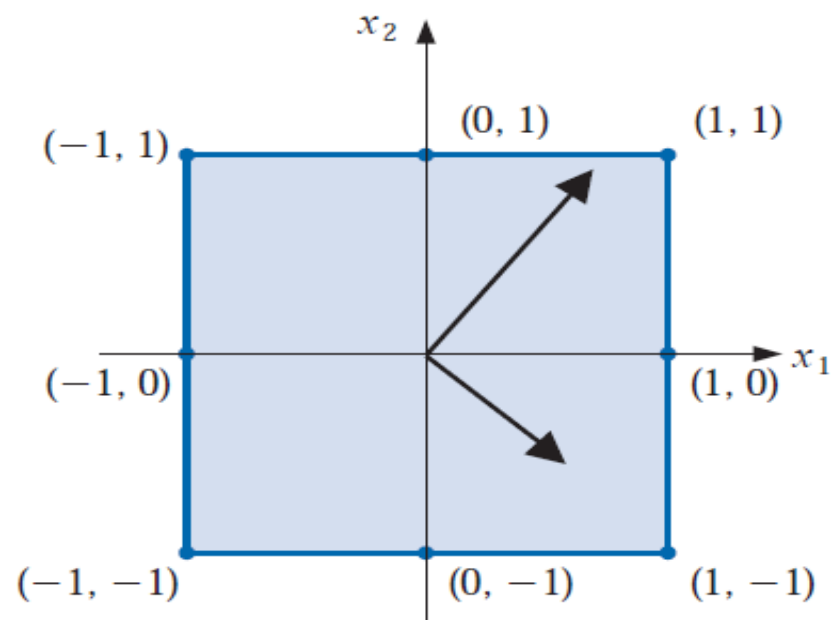
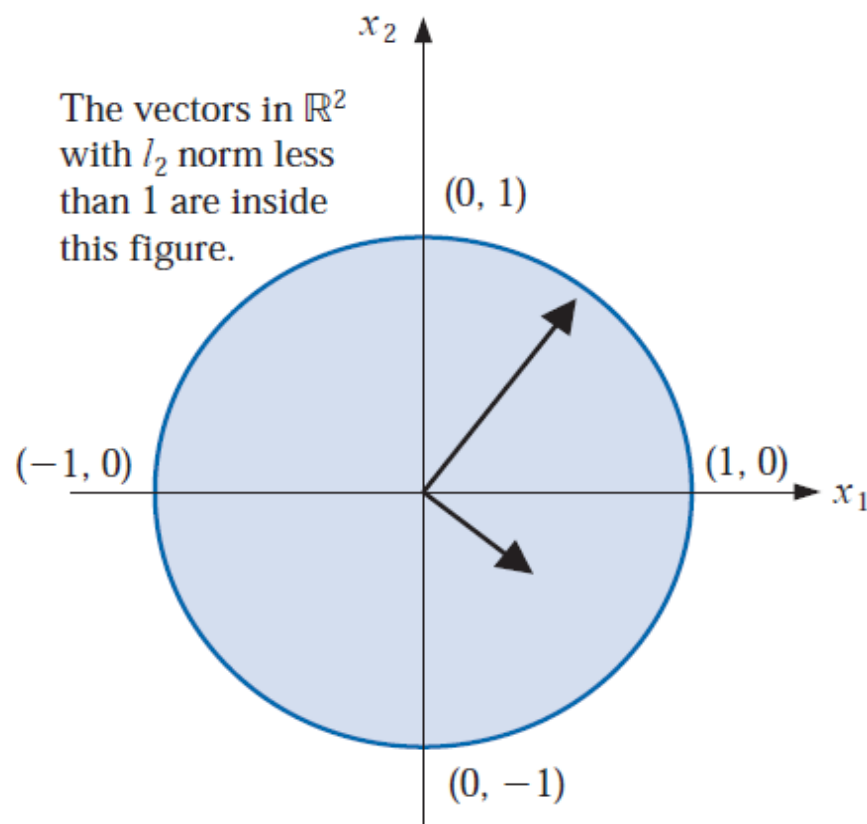
- (i)  $\| \mathbf{x} \| \geq 0$  for all  $\mathbf{x} \in \mathbb{R}^n$ ,
- (ii)  $\| \mathbf{x} \| = 0$  if and only if  $\mathbf{x} = \mathbf{0}$ ,
- (iii)  $\| \alpha \mathbf{x} \| = |\alpha| \| \mathbf{x} \|$  for all  $\alpha \in \mathbb{R}$  and  $\mathbf{x} \in \mathbb{R}^n$ ,
- (iv)  $\| \mathbf{x} + \mathbf{y} \| \leq \| \mathbf{x} \| + \| \mathbf{y} \|$  for all  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^n$ .

## Definition

The  $l_2$  and  $l_\infty$  norms for the vector  $\mathbf{x} = (x_1, x_2, \dots, x_n)^t$  are defined by

$$\|\mathbf{x}\|_2 = \left\{ \sum_{i=1}^n x_i^2 \right\}^{1/2} \quad \text{and} \quad \|\mathbf{x}\|_\infty = \max_{1 \leq i \leq n} |x_i|$$

The  $l_2$  norm is called the **Euclidean norm** of the vector  $\mathbf{x}$



The vectors in  $\mathbb{R}^2$  with  $l_\infty$  norm less than 1 are inside this figure.

## Example

Determine the  $l_2$  norm and the  $l_\infty$  norm of the vector  $\mathbf{x} = (-1, 1, -2)^t$ .

**Solution** The vector  $\mathbf{x} = (-1, 1, -2)^t$  in  $\mathbb{R}^3$  has norms

$$\|\mathbf{x}\|_2 = \sqrt{(-1)^2 + (1)^2 + (-2)^2} = \sqrt{6}$$

and

$$\|\mathbf{x}\|_\infty = \max\{|-1|, |1|, |-2|\} = 2.$$

## Distance between Vectors in $\mathbb{R}^n$

If  $\mathbf{x} = (x_1, x_2, \dots, x_n)^t$  and  $\mathbf{y} = (y_1, y_2, \dots, y_n)^t$  are vectors in  $\mathbb{R}^n$ , the  $l_2$  and  $l_\infty$  distances between  $\mathbf{x}$  and  $\mathbf{y}$  are defined by

$$\|\mathbf{x} - \mathbf{y}\|_2 = \left\{ \sum_{i=1}^n (x_i - y_i)^2 \right\}^{1/2} \quad \text{and} \quad \|\mathbf{x} - \mathbf{y}\|_\infty = \max_{1 \leq i \leq n} |x_i - y_i|$$



## Definition

A sequence  $\{\mathbf{x}^{(k)}\}_{k=1}^{\infty}$  of vectors in  $\mathbb{R}^n$  is said to **converge** to  $\mathbf{x}$  with respect to the norm  $\|\cdot\|$  if, given any  $\varepsilon > 0$ , there exists an integer  $N(\varepsilon)$  such that

$$\|\mathbf{x}^{(k)} - \mathbf{x}\| < \varepsilon, \quad \text{for all } k \geq N(\varepsilon)$$

## Jacobi's Method

The **Jacobi iterative method** is obtained by solving the  $i$ th equation in  $\mathbf{Ax} = \mathbf{b}$  for  $x_i$  to obtain (provided  $a_{ii} \neq 0$ )

$$x_i^{(k)} = \frac{1}{a_{ii}} \left[ \sum_{\substack{j=1 \\ j \neq i}}^n \left( -a_{ij} x_j^{(k-1)} \right) + b_i \right], \quad \text{for } i = 1, 2, \dots, n$$

$\mathbf{x}^{(0)}$  is an initial approximation to  $\mathbf{x}$ .

## Example

The linear system  $A\mathbf{x} = \mathbf{b}$  given by

$$E_1 : \quad 10x_1 - x_2 + 2x_3 = 6,$$

$$E_2 : \quad -x_1 + 11x_2 - x_3 + 3x_4 = 25,$$

$$E_3 : \quad 2x_1 - x_2 + 10x_3 - x_4 = -11,$$

$$E_4 : \quad 3x_2 - x_3 + 8x_4 = 15$$

has the unique solution  $\mathbf{x} = (1, 2, -1, 1)^t$ . Use Jacobi's iterative technique to find approximations  $\mathbf{x}^{(k)}$  to  $\mathbf{x}$  starting with  $\mathbf{x}^{(0)} = (0, 0, 0, 0)^t$  until

$$\frac{\|\mathbf{x}^{(k)} - \mathbf{x}^{(k-1)}\|_\infty}{\|\mathbf{x}^{(k)}\|_\infty} < 10^{-3}$$

**Solution** We first solve equation  $E_i$  for  $x_i$ , for each  $i = 1, 2, 3, 4$ , to obtain

$$\begin{aligned}x_1 &= \frac{1}{10}x_2 - \frac{1}{5}x_3 + \frac{3}{5}, \\x_2 &= \frac{1}{11}x_1 + \frac{1}{11}x_3 - \frac{3}{11}x_4 + \frac{25}{11}, \\x_3 &= -\frac{1}{5}x_1 + \frac{1}{10}x_2 + \frac{1}{10}x_4 - \frac{11}{10}, \\x_4 &= -\frac{3}{8}x_2 + \frac{1}{8}x_3 + \frac{15}{8}.\end{aligned}$$

From the initial approximation  $\mathbf{x}^{(0)} = (0, 0, 0, 0)^t$  we have  $\mathbf{x}^{(1)}$  given by

$$\begin{aligned}x_1^{(1)} &= \frac{1}{10}x_2^{(0)} - \frac{1}{5}x_3^{(0)} + \frac{3}{5} = 0.6000, \\x_2^{(1)} &= \frac{1}{11}x_1^{(0)} + \frac{1}{11}x_3^{(0)} - \frac{3}{11}x_4^{(0)} + \frac{25}{11} = 2.2727, \\x_3^{(1)} &= -\frac{1}{5}x_1^{(0)} + \frac{1}{10}x_2^{(0)} + \frac{1}{10}x_4^{(0)} - \frac{11}{10} = -1.1000, \\x_4^{(1)} &= -\frac{3}{8}x_2^{(0)} + \frac{1}{8}x_3^{(0)} + \frac{15}{8} = 1.8750.\end{aligned}$$

Additional iterates,  $\mathbf{x}^{(k)} = (x_1^{(k)}, x_2^{(k)}, x_3^{(k)}, x_4^{(k)})^t$ , are presented in the below table:

$k$	0	1	2	3	4	5	6	7	8	9	10
$x_1^{(k)}$	0.0000	0.6000	1.0473	0.9326	1.0152	0.9890	1.0032	0.9981	1.0006	0.9997	1.0001
$x_2^{(k)}$	0.0000	2.2727	1.7159	2.053	1.9537	2.0114	1.9922	2.0023	1.9987	2.0004	1.9998
$x_3^{(k)}$	0.0000	-1.1000	-0.8052	-1.0493	-0.9681	-1.0103	-0.9945	-1.0020	-0.9990	-1.0004	-0.9998
$x_4^{(k)}$	0.0000	1.8750	0.8852	1.1309	0.9739	1.0214	0.9944	1.0036	0.9989	1.0006	0.9998

We stopped after ten iterations because

$$\frac{\|\mathbf{x}^{(10)} - \mathbf{x}^{(9)}\|_{\infty}}{\|\mathbf{x}^{(10)}\|_{\infty}} = \frac{8.0 \times 10^{-4}}{1.9998} < 10^{-3}$$

The Jacobi method can be written as:

$$\mathbf{x}^{(k)} = T\mathbf{x}^{(k-1)} + \mathbf{c}$$

In the previous example,

$$\begin{aligned}x_1 &= \frac{1}{10}x_2 - \frac{1}{5}x_3 + \frac{3}{5}, \\x_2 &= \frac{1}{11}x_1 + \frac{1}{11}x_3 - \frac{3}{11}x_4 + \frac{25}{11}, \\x_3 &= -\frac{1}{5}x_1 + \frac{1}{10}x_2 + \frac{1}{10}x_4 - \frac{11}{10}, \\x_4 &= -\frac{3}{8}x_2 + \frac{1}{8}x_3 + \frac{15}{8}.\end{aligned}$$

$$T = \begin{bmatrix} 0 & \frac{1}{10} & -\frac{1}{5} & 0 \\ \frac{1}{11} & 0 & \frac{1}{11} & -\frac{3}{11} \\ -\frac{1}{5} & \frac{1}{10} & 0 & \frac{1}{10} \\ 0 & -\frac{3}{8} & \frac{1}{8} & 0 \end{bmatrix} \quad \text{and} \quad \mathbf{c} = \begin{bmatrix} \frac{3}{5} \\ \frac{25}{11} \\ -\frac{11}{10} \\ \frac{15}{8} \end{bmatrix}$$

Suppose that,

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}$$

$$A = \begin{bmatrix} a_{11} & 0 & \cdots & 0 \\ 0 & a_{22} & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & a_{nn} \end{bmatrix} - \begin{bmatrix} 0 & \cdots & \cdots & 0 \\ -a_{21} & \cdots & \cdots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ -a_{n1} & \cdots & -a_{n,n-1} & 0 \end{bmatrix} - \begin{bmatrix} 0 & -a_{12} & \cdots & -a_{1n} \\ \vdots & \ddots & \ddots & \vdots \\ \vdots & \cdots & -a_{n-1,n} & \vdots \\ 0 & \cdots & \cdots & 0 \end{bmatrix}$$

$$= D - L - U.$$

$$A\mathbf{x} = \mathbf{b} \implies (D - L - U)\mathbf{x} = \mathbf{b} \implies D\mathbf{x} = (L + U)\mathbf{x} + \mathbf{b}$$

and, if  $D^{-1}$  exists, that is, if  $a_{ii} \neq 0$  for each  $i$ , then

$$\mathbf{x} = D^{-1}(L + U)\mathbf{x} + D^{-1}\mathbf{b}$$

This results in the matrix form of the Jacobi iterative technique:

$$\mathbf{x}^{(k)} = D^{-1}(L + U)\mathbf{x}^{(k-1)} + D^{-1}\mathbf{b}, \quad k = 1, 2, \dots$$

Introducing the notation  $T_j = D^{-1}(L + U)$  and  $\mathbf{c}_j = D^{-1}\mathbf{b}$  gives

$$\mathbf{x}^{(k)} = T_j\mathbf{x}^{(k-1)} + \mathbf{c}_j$$

## The Gauss-Seidel Method

$$A\mathbf{x} = \mathbf{b}$$

$$x_i^{(k)} = \frac{1}{a_{ii}} \left[ - \sum_{j=1}^{i-1} (a_{ij}x_j^{(k)}) - \sum_{j=i+1}^n (a_{ij}x_j^{(k-1)}) + b_i \right] \quad i = 1, 2, \dots, n$$